

# Outline

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| FOL                                   | TPTP                               |
|---------------------------------------|------------------------------------|
| $\perp, \top$                         | <code>\$false, \$true</code>       |
| $\neg F$                              | <code>~F</code>                    |
| $F_1 \wedge \dots \wedge F_n$         | <code>F1 &amp; ... &amp; Fn</code> |
| $F_1 \vee \dots \vee F_n$             | <code>F1   ...   Fn</code>         |
| $F_1 \rightarrow F_n$                 | <code>F1 =&gt; Fn</code>           |
| $(\forall x_1) \dots (\forall x_n) F$ | <code>! [X1, ..., Xn] : F</code>   |
| $(\exists x_1) \dots (\exists x_n) F$ | <code>? [X1, ..., Xn] : F</code>   |

# More on the TPTP Syntax

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%---- 1 * x = 1
fof(left_identity,axiom,(
    ! [X] : mult(e,X) = X )).
%---- i(x) * x = 1
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%---- (x * y) * z = x * (y * z)
fof(associativity,axiom,(
    ! [X,Y,Z] :
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%---- x * x = 1
fof(group_of_order_2,hypothesis,
    ! [X] : mult(X,X) = e ).
%---- prove x * y = y * x
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202. sP1(mult(sK,sK0)) [backward demodulation 188,15]
188. mult(X8,X9) = mult(X9,X8) [superposition 22,87]
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14. ~sP1(mult(sK,sK0)) [inequality splitting name introduction]
13. mult(sK,sK0) != mult(sK0,sK) [cnf transformation 8]
12. e = mult(X0,X0) (0:5) [cnf transformation 4]
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- ▶ Each inference derives a formula from zero or more other formulas;
- ▶ Input, preprocessing, new symbols introduction, superposition calculus
- ▶ **Proof by refutation**, generating and simplifying inferences, unused formulas ...

# Proof by Vampire (Slightly Modified)

Refutation found. Thanks to Tanya!

```
203. $false [subsumption resolution 202,14]
202. sP1(mult(sK,sK0)) [backward demodulation 188,15]
188. mult(X8,X9) = mult(X9,X8) [superposition 22,87]
87. mult(X2,mult(X1,X2)) = X1 [forward demodulation 71,27]
71. mult(inverse(X1),e) = mult(X2,mult(X1,X2)) [superposition 23,20]
27. mult(inverse(X2),e) = X2 [superposition 22,10]
23. mult(inverse(X4),mult(X4,X5)) = X5 [forward demodulation 18,9]
22. mult(X0,mult(X0,X1)) = X1 [forward demodulation 16,9]
20. e = mult(X0,mult(X1,mult(X0,X1))) [superposition 11,12]
18. mult(e,X5) = mult(inverse(X4),mult(X4,X5)) [superposition 11,10]
16. mult(e,X1) = mult(X0,mult(X0,X1)) [superposition 11,12]
15. sP1(mult(sK0,sK)) [inequality splitting 13,14]
14. ~sP1(mult(sK,sK0)) [inequality splitting name introduction]
13. mult(sK,sK0) != mult(sK0,sK) [cnf transformation 8]
12. e = mult(X0,X0) (0:5) [cnf transformation 4]
11. mult(mult(X0,X1),X2)=mult(X0,mult(X1,X2)) [cnf transformation 3]
10. e = mult(inverse(X0),X0) [cnf transformation 2]
9. mult(e,X0) = X0 [cnf transformation 1]
8. mult(sK,sK0) != mult(sK0,sK) [skolemisation 7]
7. ? [X0,X1] : mult(X0,X1) != mult(X1,X0) [ennf transformation 6]
6. ~! [X0,X1] : mult(X0,X1) = mult(X1,X0) [negated conjecture 5]
5. ! [X0,X1] : mult(X0,X1) = mult(X1,X0) [input]
4. ! [X0] : e = mult(X0,X0) [input]
3. ! [X0,X1,X2] : mult(mult(X0,X1),X2) = mult(X0,mult(X1,X2)) [input]
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# Statistics

Version: Vampire 1.8 (revision 1362)  
Termination reason: Refutation

Active clauses: 14  
Passive clauses: 35  
Generated clauses: 194  
Final active clauses: 8  
Final passive clauses: 11  
Input formulas: 5  
Initial clauses: 6

Splitted inequalities: 1

Fw subsumption resolutions: 1  
Fw demodulations: 68  
Bw demodulations: 14

Forward subsumptions: 65  
Backward subsumptions: 1  
Fw demodulations to eq. taut.: 20  
Bw demodulations to eq. taut.: 1

Forward superposition: 60  
Backward superposition: 39  
Self superposition: 6

Unique components: 6

Memory used [KB]: 255  
Time elapsed: 0.007 s

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- ▶ Chris Weidenbach: **A dark side of theorem proving.**
- ▶ Anonymous referee: **VAMPIRE is not a glass of Tuborg.**

# Main applications

- ▶ Software and hardware verification;
- ▶ Static analysis of programs;
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- ▶ Reasoning in non-classical logics;
- ▶ Program synthesis;
- ▶ Writing papers and giving talks at various conferences and schools ...

# What an Automatic Theorem Prover is Expected to Do

## Input:

- ▶ a set of **axioms** (first order formulas) or clauses;
- ▶ a **conjecture** (first-order formula or set of clauses).

## Output:

- ▶ **proof** (hopefully).

# Proof by Refutation

Given a problem with axioms and assumptions  $F_1, \dots, F_n$  and conjecture  $G$ ,

1. negate the conjecture;
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Thus, we reduce the theorem proving problem to the problem of **checking unsatisfiability**.

In this formulation the negation of the conjecture  $\neg G$  is treated like any other formula. In fact, Vampire (and other provers) **internally treat conjectures differently, to make proof search more goal-oriented**.

# General Scheme (simplified)

- ▶ Read a problem;
- ▶ Determine proof-search options to be used for this problem;
- ▶ Preprocess the problem;
- ▶ Convert it into CNF;
- ▶ Run a saturation algorithm on it, try to derive  $\perp$ .
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Trying to derive  $\perp$  using a saturation algorithm is the hardest part, which in practice may not terminate or run out of memory.